Each problem is worth 10 points. Budget your time carefully, the last two problems may take more time than the first two problems.

- 1. A. Evaluate the following simple Lisp Expressions:
- (i) (cons '(a) '(b))

- (ii) (append '(a) '(b))
- (iii) (car '((a b) c d))
- (iv) (atom '(a))
- (v) (cond (nil "A") (t "B"))
- B. Draw the binary tree representation for x: (setq x '(((a b) c (d)) e)) and figure out the correct car cdr sequence to obtain b from x.

- 2. Label the statements below as to which principle best fits. (Choose from Abstraction, Automation, Labeling, Orthogonality, Regularity and Security.)
- A. Make the computer/compiler/interpreter do the tedious chores.
- B. Factor out the recurring patterns.
- C. For example, indexing was used to illustrate this principle in Chapter 1.
- D. Independent functions should be controlled by independent mechanisms.
- E. Leads to the notion of binding a symbolic name to a location.

```
3. Consider the Lisp functions A and B below:
(def A (lambda (x)
                                      ;; x is a list
            (cond ((null x) nil)
                   ((null (cdr x)) t)
                   (t nil))))
(def B (lambda (x)
                                      ;; x is a list of lists
            (cond ((null x) t)
                   ((A (car x)) (B (cdr x)))
                   (t nil))))
Evaluate:
(i) (A '(c))
(ii) (A '(c d))
(iii) (A '((c)))
(iv) (B '((a)(b)))
```

(v) (B '(((a b))))

4. Write a recursive Lisp boolean function all-atoms with one parameter x which is a list, the function returns true if x is nil or a list of atoms.

Quiz	2
COP	4020

by				
-	6	Oct	RO	

Each problem is worth 10 points. Budget your time carefully, the quiz is on the long side.

1 Terminology: Match the letter of the phrase (below) that <u>best</u> applies to the following terms:

associativity ____ aliasing ____ binding ___ call by reference ___ call by value ____ key word

____ overloading ____ precedence scope of a binding

- A. A word that has special meaning in certain syntactic contexts.
- B. A word that can't be used for a programmer-declared object.
- C. Rule for grouping the same operator.
- D. Rule for priority of operators.
- E. A single data object with two or more names.
- F. A single name denoting more than one thing within the same scope.
- G. Associates names with locations or objects.
- H. That region of the program in which it is visible.
- I. Actual parameters are copied into the formal parameters.
- J. Addresses of actual parameters are used for formal parameters.
- 2. For the C code in the middle:
- A. Draw a contour diagram (show parameters and procedures names too).

```
main()
{
  int x = 5;
  P(7);
}
P(int y)
{
  int z = 2;
  Q(11);
}
Q(int w)
{
  /*show*/
}
```

reserve word

```
3. Grammar: Consider the rules
                                    S := A | L | (S, S)
                                    L ::= (S) | (S L)
For the strings below either
give a deivation tree, or state
                                    A := a | b | c | d
that it cannot be devired. (Here S is the start symbol and A & L are
non-terminals.)
         B. (a, (b, (c, d))) C. (a ((b, (c)))) D. (a (b, c))
A. (a b)
4. Lisp
A. If (def partA (lambda (x)
           (cond ((atom x) x)
                  (t (cons (partA (car x)) (partA (cdr x))))))
evaluate (partA '(a ((b) c) d) )
B. Here (<= n m) is true if and only if "n <= m"
If (def partB (lambda (x y)
            (cond ((null x) y)
                  ((null y) x)
                  ((<= (car x) (car y)) (cons (car x) (partB (cdr x) y))
```

(t (cons (car y) (partB x (cdr y))))))

evaluate (partB '(1 5 7) '(4 9))

Quiz 3	by	
COP 4020		20 Oct 89
Each problem is worth quiz (as always) is on		your time carefully, the
1 Fill in the blanks:		
A. A coercion is a type co		version and a cast is a
B. Algol used two form	ms of parameter pass	sing, call by value and call
by In C, b	y the use of macros	(#define's with parameters)
		s passed by
C. When are two typ	es identical? The m	ost restrictive definition is hile the least restrictive
definition is called _	equivalence, w	uivalence
D. Scoping rules dete	ermine in which envir	onment are non-local
identifiers accessed.	If it is in the enviro	nment of definition, it is
called	scoping and if it is	in the environment of its
caller, it is called	scoping) .
E. In C, Pascal and A	llgol, local variables a	re bound to absolute
		oal variables are bound to
absolute addresses a	.T Tim	e.
2. Grammar: Consider	the rules S ::= A	(T .
A. For the strings bel	•	•

A ::= a | b | c | d

(ii). (c (d a) b) (i). (a b)

give a derivation tree, or state

non-terminals.)

B. Show the grammar is ambiguous. (Find two different derivation trees for some string in the language.)

that it cannot be devired. (Here S is the start symbol and A & T are

Program Contour; Var a, b:integer;
Procedure One(var n, m: integer); Var c: integer;
<pre>Procedure Xyz(var y, z: integer); begin ***here*** end;</pre>
Procedure Abc(var r, s: integer); begin Xyx(r, s) end;
begin Abc(c, m) end;
Procedure Two(var i, j:integer); Var d:integer; begin One(i, d) end;
begin Two(a, b) end.
3. Draw a contour diagram for the Pascal-like code above (include parameter and procedure names too).
4. Aliases: When execution reaches the line marked ***here*** in the Pascal-like code above each procedure has a activation record. Give all the aliases for the each of the variables (even if they are out of scope at this ***here*** line). Note if w "is an alias" for x, then x "is an alias" for w. a: b:
a: b: c: d:
i: j:
m: n: s:
y: z:

Quiz 4	by		
COP 4020			13 Nov 89
	worth 10 points. Bus on the long side.	idget your time caref	ully, the
applies to the fo mutual conse derived type package		structural e	alence ameters
type, but is disti B. Has the same subset of the don C. When used in D. Two objects a with the same ty E. Two objects ar from the same do F. Actual parame G. Actual parame H. Can be implem I. When used in	nct from it. attributes and operanain. Ada it creates aliase re of the same type re identifer. The of the same type omain of values. The ters can be listed in the same and the defendance and the same and the defendance and the same and the	if they are declared if they are declared any order. ate procedure.	but is a together or as coming
parameters or	eter modes are parameted param	parameters or ters. allow parameters which are with the determined by exception ral instances of a stack of intergers and ations like push and	ch are functions. statement, (for

W - 1

Procedure main is declare c, d: float; Package A is Function X (c, d: in float) return float is begin **INNER** ... end; Function Y (p: in float := 0.0) return float is begin ... end; End A; Package B is Function X (c, d: in integer := 1) return integer is begin ... end; Function Y (p: in integer := 0) return integer is begin ... end; End B; begin ***HERE*** end main: 3. For the Ada like code above A. (4pts) At the ***HERE*** which function would be called by (ii) Y (p=>4.5)(i) Y (45)_____ (iv) Y () (iii) X (c,d) B. (2pts) At the **INNER** is it possible to "see" the variable c in main, if so write a statement which will assign to c in main the value of the parameter d in A.X.

- C. (4pts) Write down four different ways in Ada to call function "+" (left, right: in colors) return colors; with red the actual parameter for left and green the actual parameter for right.
- 4. Project: Use C to write a recursive-descent recognizer for the grammer below. Assume token is the next CHARACTER in the input stream. Assume the function advance(); advances the token to the next character. Assume a main() which has already called advance() once. Assume the input stream has no white space or newline characters. Write the boolean functions have U (and respectively, have V) which return true or false depending on if the input is a string in U (respectively, in V). (Sort of like get_s and get_t but they return true or false instead of anything useful.) U is the start symbol, and V is the only other nonterminal.

U ::= VV | @ V ::= U\$ | #

Quiz 5	by		
COP 4020		1 Dec 89	
Each problem i #include "stan	s worth 10 points. Budg dard_disclaimer_on_len	jet your time carefully. igth_of_quiz"	
applies to the	following terms:	phrase (below) that best	
class	extent side effect programming language	<pre>pure function garbage collection dangling reference</pre>	
effects (refere B. One whose C. That region	ntial transparency). output is solely determin of the program in which	functions, so there are no side- ned by its input. a binding is visible. the storage used to hold a	
F. An access p exceeds extent G. A pool of st	non-local variable by a ath to a data structure the control of the	function or a procedure. hat has been destroyed (scope are created and destroyed in	S.
•		are created and destroyed in	
I. When used in J. Automatic a	n C++ it can create an al nd costly recovery of st	bstract data type. orage from data objects which g out all possible reference	
2. Project (Lisp A. (pairlis '(w x	o): evaluate: y z) '(a 7 (c) nil))		
B. (assoc 's '((t.7)(u lambda (x) (y))(w	.s)(s.5)(t.s)(s.a)(good.doctor)))	
Ç. (mapcar '(la	mbda (x) (times x (minus	; x 1))) '(2 4 6 8))	
D. (label fn (la	mbda (x) ero x) 0)(t (plus x (fn (m	ninus x 1))))) 5)	

3. Fill in the blanks: A. In Ada programming is done by use of tasks. In a task a select statement allows the task to execute any of its which another task has called. When two tasks interact it is called a
B. In C++ the objects are called, while these are similar to a structure with a typedef but they can also have functions and procedures which are also called in object oriented lingo. C. Exceptions in Ada can be started with the statement, the code which is then executed is determined by scoping, and this code is called the exception D. Information hiding or encapsulation is the practice of hiding information that the user doesn't need to know about the implementation of the abstraction. Two reasons in favor of this are:
4. A. The DoD is moving towards requiring Ada as its only language. Give 3 reasons in favor of such a move.

B. Give 3 reasons against such a move.

Mid Term COP 4020	by	27 Oct 89
Each problem is worth 10 p	ooints. Budget your tim	e carefully.
1 Terminology: Match the leapplies to the following terminology: Match the leapplies to the following terminology: call by name overloading side effect	ms: aliasing name	equivalence erminal symbol
A. The lefthand side of a BN B. Associates locations with C. Has the same attributes a subset of the domain. D. Two objects are of the swith the same type identified E. A single data object with F. A single name denoting in G. Two objects are of the safetiment of the same domain of value of the same domain of value. Syntactic ambiguity issued identifiers left to be bound J. All applied occurrences of the expression of the corresponded identifiers bound	and operations as some ame type if they are deer. two or more names. The they are dealers are than one thing with ame type if they are dealers. The formal parameter are the formal parameter, with a in the local environment of the formal parameter are sponding actual parameter.	eclared together or nin the same scope. eclared as coming ments. are replaced by the ny embedded ent. are replaced by eter, with any
2. Lisp 1: Draw binary tree expressions: (a (b c)) (a (b . c)) (a (b	·	-

et .

3. Fill in the blanks:	
A. A is an im	plicit type conversion and a
is an explicit type conversio	on.
B. Two methods of impleme	enting static scoping for Pascal were
given in the text, one used	static chaining and the other was called
	er method is needed in C since C does
not allow nested	<u>-</u>
	egin-end pairs in Pascal and Algol: in
	used just to delimit
	they also delimit
	ne to a location is called a,
	econd name to the same location is
called a	I I I I I I I I I I I I I I I I I I I
E. In C, Pascal and Algol, lo	cal variables are bound to absolute
	me, while global variables are bound to
absolute addresses at	time.
1 For the Coods in the mid	410.
4. For the C code in the mide	$\int \ln t n = 35$; B. Draw the stack showing
A. Draw a contour diagram	main() the activation records,
(show parameters and	$\begin{cases} \text{int } x = 5; \text{ parameters at the "show"} \end{cases}$
procedures names too).	$P^{(7)}; \qquad \text{line. (Use "RA back to } \underline{\hspace{1cm}}$
	for return addresses \
	IP(int y) Hor return addresses.)
	int $z = 2$;
	Q(11);
	Q(int w)
	{
	/*show*/
	j

.

```
Program MidTerm; Var i, j:integer; a: array[1..2] of integer;
Procedure One; begin i := 2; end;
Procedure Two; Var i:integer; begin One; end;
Procedure Swap ( x, y:integer ); Var i, t: integer;
begin t := x; x := y; y := t; end.
begin Two; a[1]:= 3; a[2]:= 1; j:= 2; Swap(j, A[j]);end.
```

- 5. For the Pascal-like code above:
- A. (i) If Dynamic scoping is used when procedure two calls procedure one, which "i" is assigned the value two?
- (ii) If Static scoping is used when procedure two calls procedure one, which "i" is assigned the value two?
- B. If the call to Swap is call by reference, what are the values of j, a[1] and a[2] after Swap returns?
- C. If the call to Swap is call by value, what are the values of j, a[1] and a[2] after Swap returns?
- D. If the call to Swap is call by name, what are the values of j, a[1] and a[2] after Swap returns

- B. Evaluate (fun '(a nil b))
- C&D. Evaluate (fun '(a ((b) c) d))

7. C errors: Point out and correct and explain five errors in the C code below (label them (a) ... (e))

```
typedef struct { int alpha; int omega } Type_t;
full of errors()
       int x = 3; int *y; Type_t t;
                                                 Type t *pt;
      t.omega = 5;
x = t->alpha;
pt.omega = x;
y = &x*
t.omega = 5;
t.omega = 5;
                                            x = pt->alpha;
                                               y = pt;
(*pt).alpha = 3;
                            t.omega = x;
                                                  *x = y;
       y = &x;
                            *y = x;
}
а
b
С
d
е
```

8. Project: Use C to write a recursive-descent recognizer for the grammer below. Assume token is the next CHARACTER in the input stream. Assume the function advance(); advances the token to the next character. Assume the input stream has no white space or newline characters. Write the boolean functions have W (and respectively, have Y) which return true or false depending on if the input is a string in W (respectively, in Y). (Sort of like get_s and get_t but they return true or false instead of anything useful.) W is the start symbol, and Y is the only other nonterminal.

W ::= #Y# Y ::= \$W | @Y | %