INTEROFFICE MEMORANDUM

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Time Warp Measurements on Message Times

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Recently, we (Mike DiLoreto and myself) were able to "synchronize the 2 microsecond clocks" across the Mark III Hypercube. This allows a number of accurate measurements of communication time. These statistics are the first published fruit of this instrumentation. Rather than means and standard deviations we offer a large number of graphs. We feel the shape of the graphs is more telling than averages distorted by the *noise* at the very slow (large delay) end.

Definitions

We time stamp each Time Warp message at four times during its life:

- (1) As soon as it is a message. This occurs roughly just before the routine deliver is called to mail the message.
- (2) Just before it is given to Mercury. This occurs just before the Mercury routine send msg is called.
- (3) When Mercury receives a message it calls the interrupt routine ioint which gives the third time stamp. (The message then goes into the Mercury receive queue.)
- (4) The last time stamp (and logging of all the time stamps) is done in the Time Warp routine nq. If we waited any longer the message could be annihilated.

Using the four time stamps above we define the latencies below:

- A. Time Warp End to End Time: as time stamp4 time stamp1.
- B. Time Warp to Mercury Time: as time stamp₂ time stamp₁.
- C. Mercury to Mercury Time: as time stamp3 time stamp2.
- D. Mercury to Time Warp Time: as time stamp4 time stamp3.

Observe that on node messages have neither time stamp2 nor time stamp3. Only during the Mercury to Mercury time is the message really in transit. The B and D times are queueing delays, the time the message sits in some queue. However, the Time Warp end to end time is roughly the transit time of a message from the application's point of view.

The Two Runs of CTLS

We collected data on two runs of CTLS on TW111. Both were 32 node runs made on the cpc5 hypercube. Over 65,000 messages were send on each run with over 50,000 messages being off node messages. In most of the graphs these two runs are named "acks=2" and "acks=20". The term acks (or max_acks) is a parameter which stops the sending of a new (positive) message when the number of unacknowledged messages exceeds max_acks. However, the two runs differed in other ways, with acks=20 run also having plimit=50, reward=1 and penalty=2. The parameter plimit is how far into the Mercury queue Time Warp will look for a message with higher priority than the object it is about to run. The penalty parameter "slows" objects which send negative messages and the reward parameter undoes the "slowing".

These two runs seem typical of CTLS and (with one exception) yield close to the same graphs. The acks=20 run was first and it was conjectured (incorrectly) that reducing acks to 2 would decrease the length and time in Mercury's receive queues. However, the two runs are quite similar.

The times for GVT messages have been filtered out. (The "collects" can have a long (13 millisecond) Mercury to Mercury delay. Also the GVT printf was turned off. (This can slow the GVT node by about 50 milliseconds per GVT message). Also each off node application message has a short returning acknowledgement message which is not logged. Thus Mercury had about twice as many messages as logged.

The Graphs

The graphs in the appendices were all made with MicroSoft's Excel which has a limit of about 100 data points per series. Sometimes fewer data points were used to better show the shape of the curve. Thus some graphs appear in three pieces: the fastest 95% of the messages, the 95% to 99.9% range and the slowest 0.1%. Often the graphs appear in both linear scales and semi-log paper to better show its features. The programs mstat and msgstat (on the Sun) were used to filter the over 50,000 messages into histograms with near 100 data points. The Excel Macro File "Hist.MS" was used to help speed the graphing.

Appendices A, B, C and D contain graphs for the times A, B, C and D above. (For example Appendix C is about Mercury to Mercury Times.) In addition, Appendix E contains graphs on queueing delays and queue lengths of Mercury's receive queue.

Time Warp will sometimes unlink a message not at the front of Mercury's receive queue. Usually these messages are system messages but about 100 application messages in each run were processed ahead of earlier arriving messages. Hence the strange wording on some of the graphs in Appendix E.

Conclusions

- 1. The major factor in the Time Warp end to end time is the Mercury to Time Warp time, the time a message waits in the Mercury receive queue. "Most messages spent most of their time waiting on the receiving node."
- 2. The decrease of max acks from 20 to 2 does not seem to reduce either the Time Warp end to end time nor the Mercury to Time Warp time. However, it does increase the Time Warp to Mercury Time by increasing the number of messages with large Time Warp to Mercury times. Indeed, we claim the second peak near 7 milliseconds in graph B-2 is evidence that there is a similar peak in the execution times of object and Time Warp "pieces". (Roughly, the time it takes the code to return to the "Tester main loop" so it can try to send the message again.)
- 3. The application CTLS does not push Mercury to its limits. Indeed, it seems that message length rather than the number of hypercube

hops is the important factor in determining transit time. There are messages with long Mercury to Mercury delays, but they all could be behind a GVT "collection". In a GVT collection, half of the cube (16 of 32 nodes in this case) sends its reply down paths whose last hop is the channel which connects nodes one and zero. Delays on the order of a dozen milliseconds are not surprizing, even when there are no other messages in the hypercube.

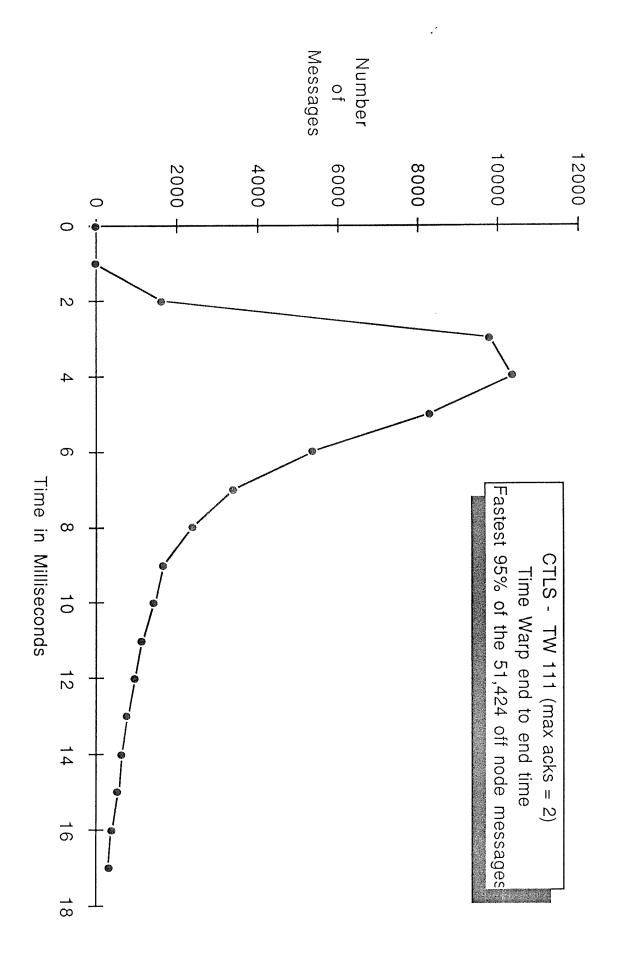
4. Although Time Warp was never designed to run in "real time," it has delays which seem to be too long. (See below, "Story of a Message.") Garbage collection on the Slooow benchmark (with one object per node and a three second GVT interval) takes 60 to 80 milliseconds. The execution time of Time Warp "pieces" clearly needs to be identified and shortened.

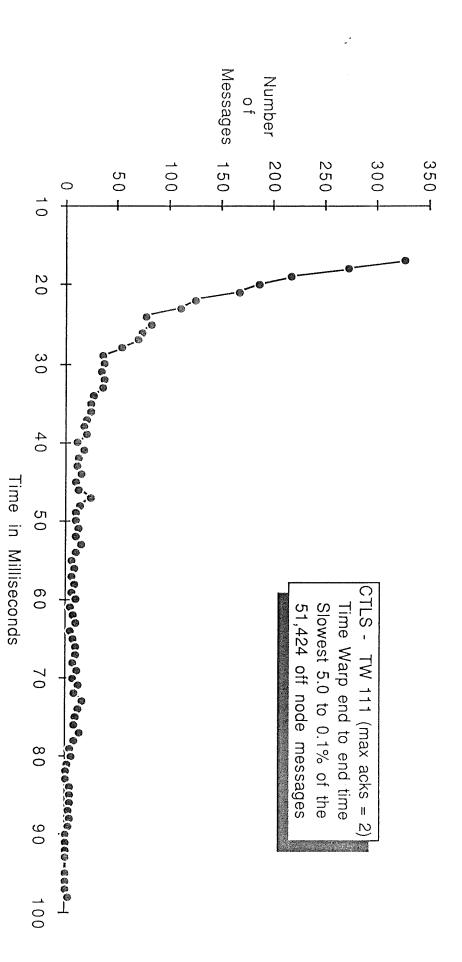
Story of a Message

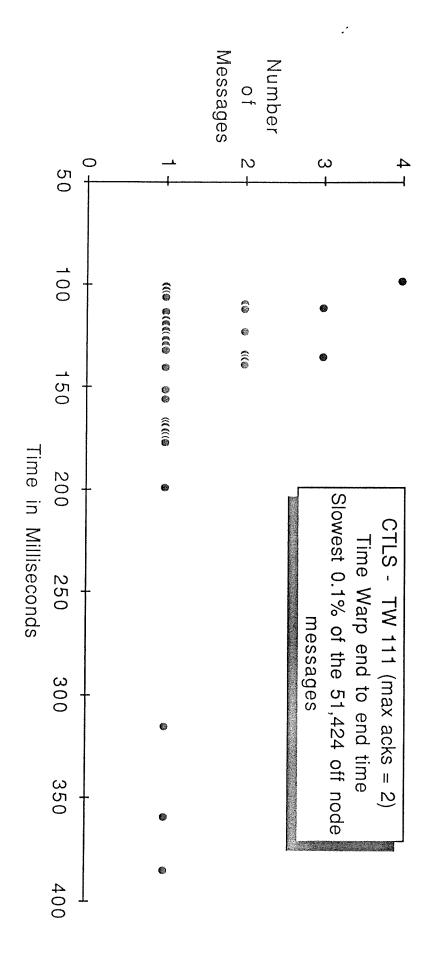
This is the history of one message which we will call "M." The message M is the right most "dot" on the graphs A-3, A-6, D-3, D-6, E-3, E-4 and E-9. M's Time Warp end to end time is 385 milliseconds of which 383 milliseconds is spent waiting in Mercury's receive queue. However, when M arrives on the receiving node there are no other messages in the Mercury receive queue. It seems that "Time Warp code" is executing during all of this time. Time Warp is busy doing "zip forward," that is it is repeating executing "go forward" for the object red_div22. The object red_div22 sends messages to itself, and lots of them are now being cancelled. There are 64 messages from red_div22 to itself which are being cancelled while message M waits in Mercury's receive queue. The irony is that M is a message for red_div22 with a receive time earlier than the send time of all but five of the 64 cancelled messages.

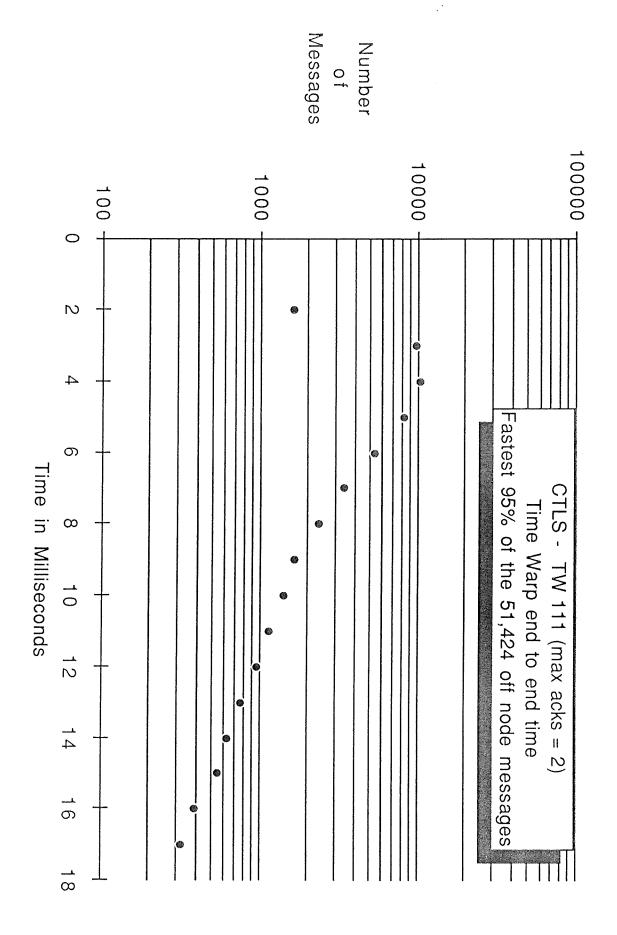
Appendix A Time Warp End to End Times

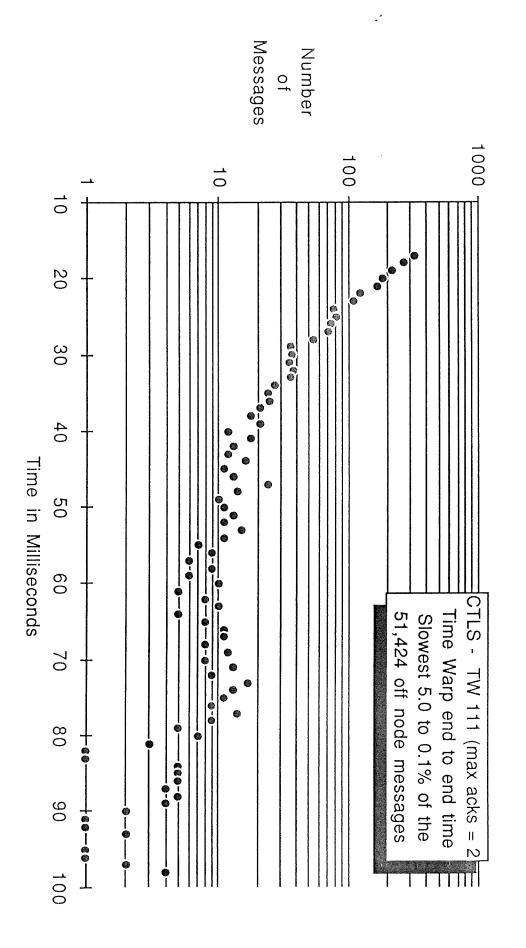
- A-1: Fastest 95%, acks=2, graph (see A-4).
- A-2: The 95% 99.9% range, acks=2, graph (see A-5).
- A-3: Slowest 0.1%, acks=2, graph (see A-6).
- A-4: Fastest 95%, acks=2, semi-log graph (see A-1).
- A-5: The 95% 99.9% range, acks=2, semi-log graph (see A-2).
- A-6: Slowest 0.1%, acks=2, semi-log graph (see A-3).
- A-7: Fastest 98%, acks=2 vs acks=20, graph (see A-8, A-9).
- A-8: Fastest 98%, acks=2 vs acks=20, semi-log graph (see A-7, A-9).
- A-9: Fastest 98%, acks=2 vs acks=20, log-log graph (see A-7, A-8).
- A-10:Two Queueing Delays? (see A-11).
- A-11:Fastest 99.9%, acks=2, graph (see A-10).

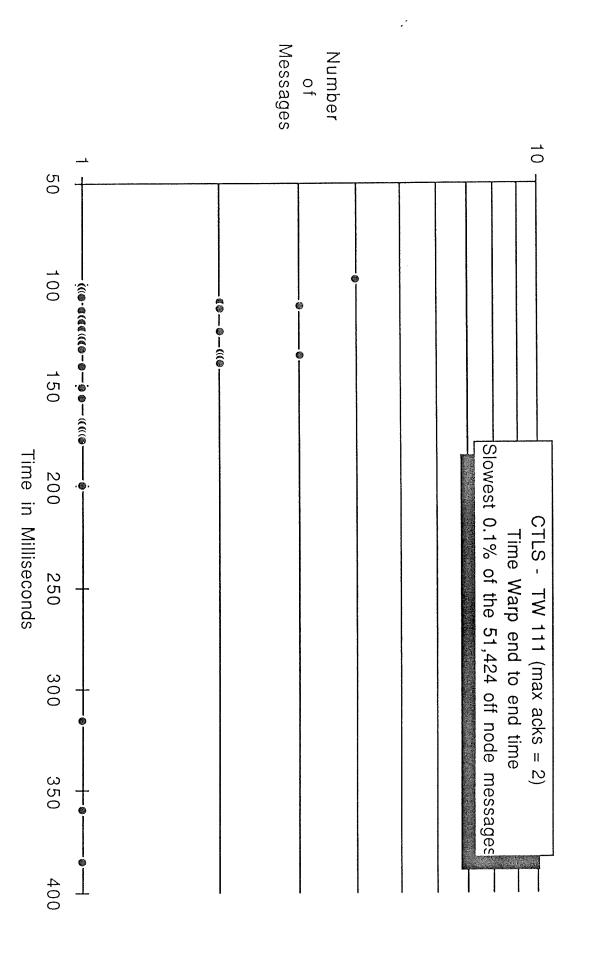


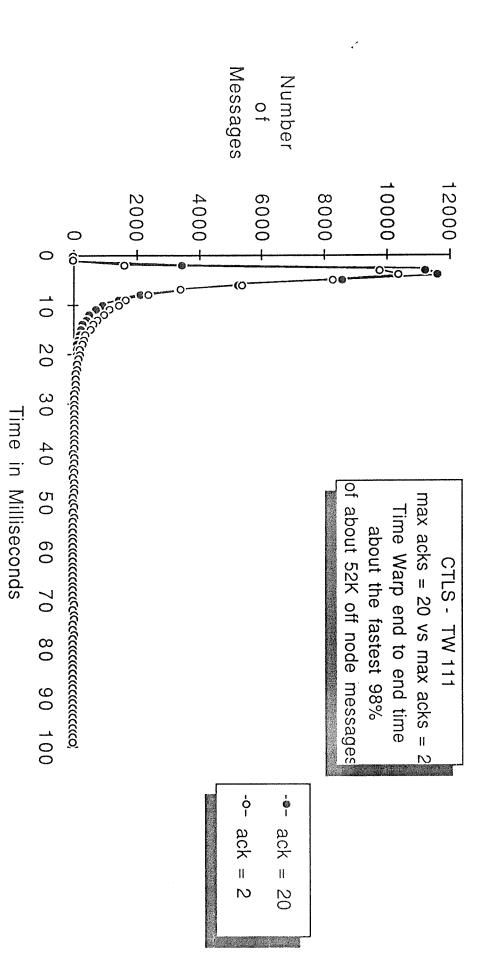


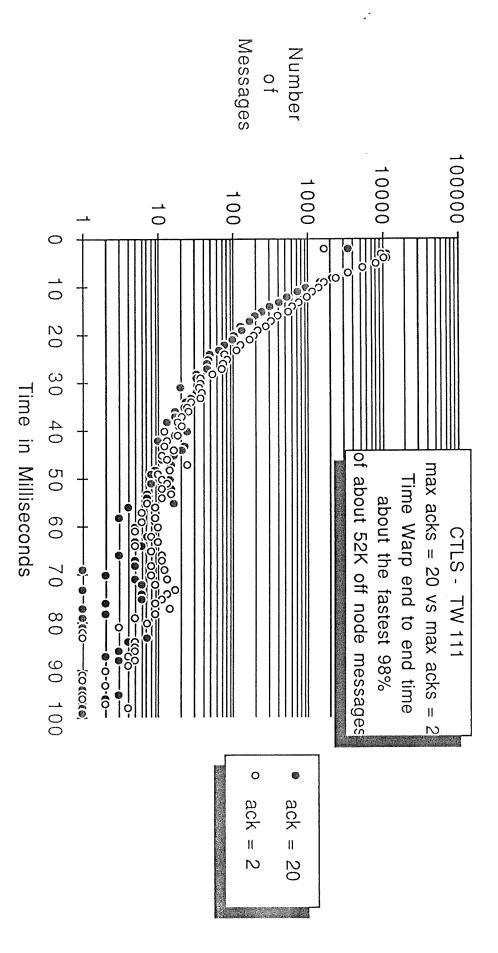


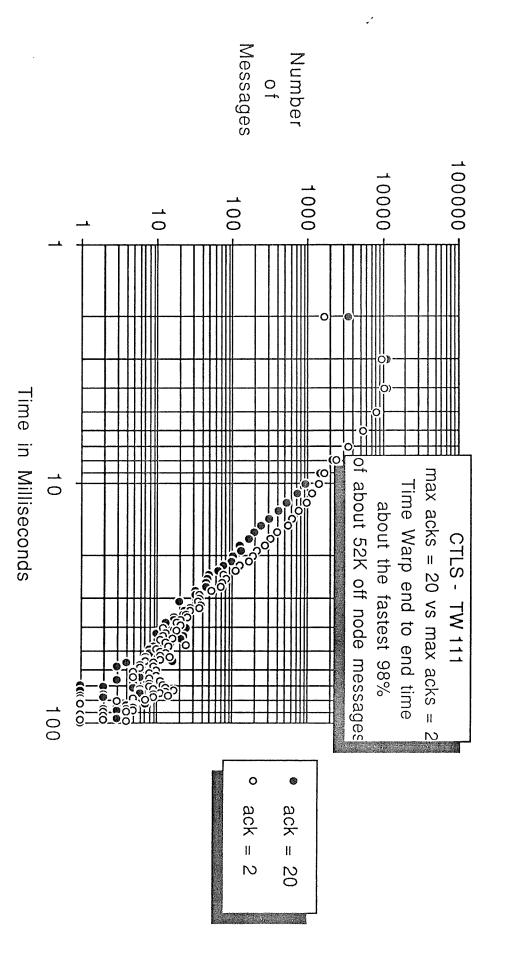




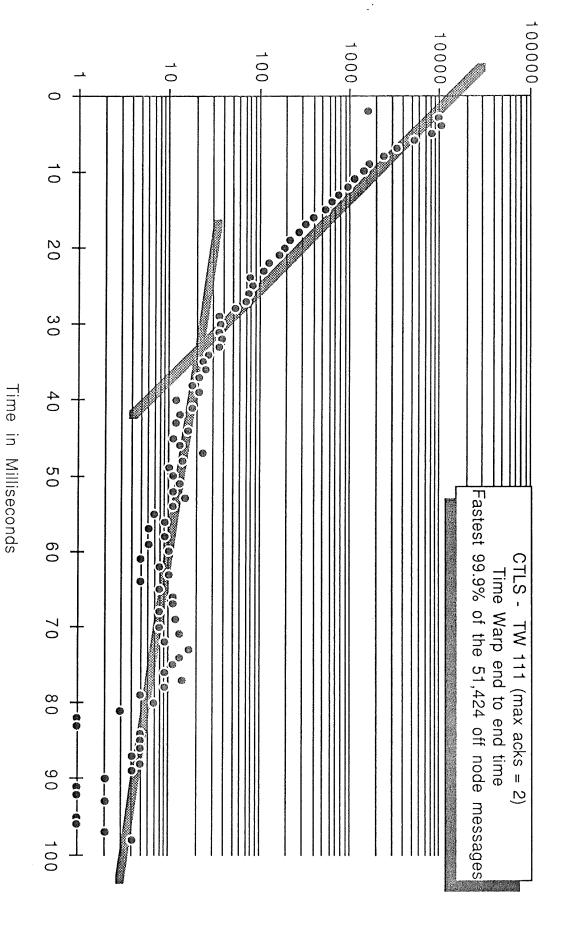


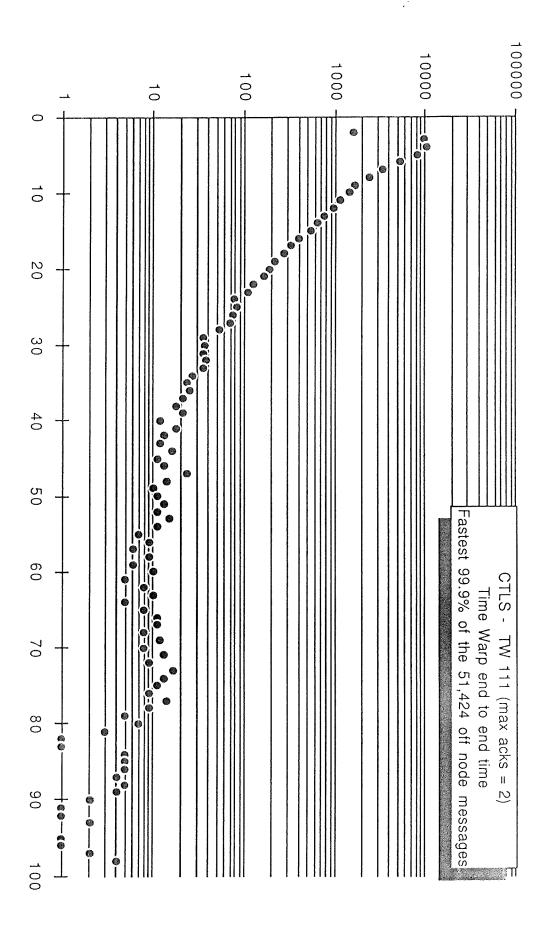






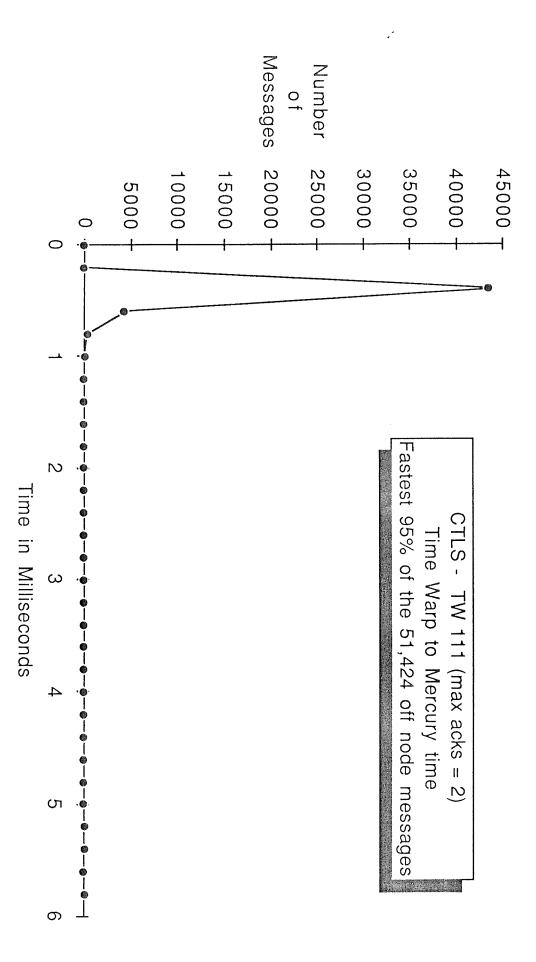
Two Queueing Delays?

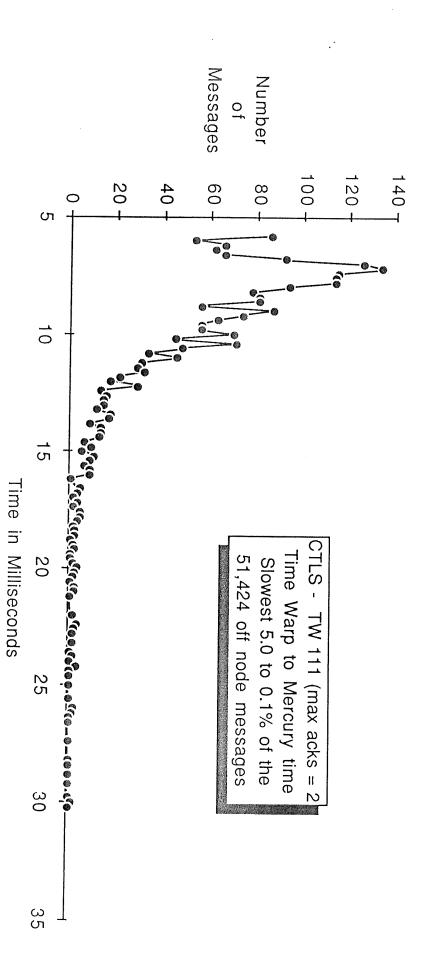


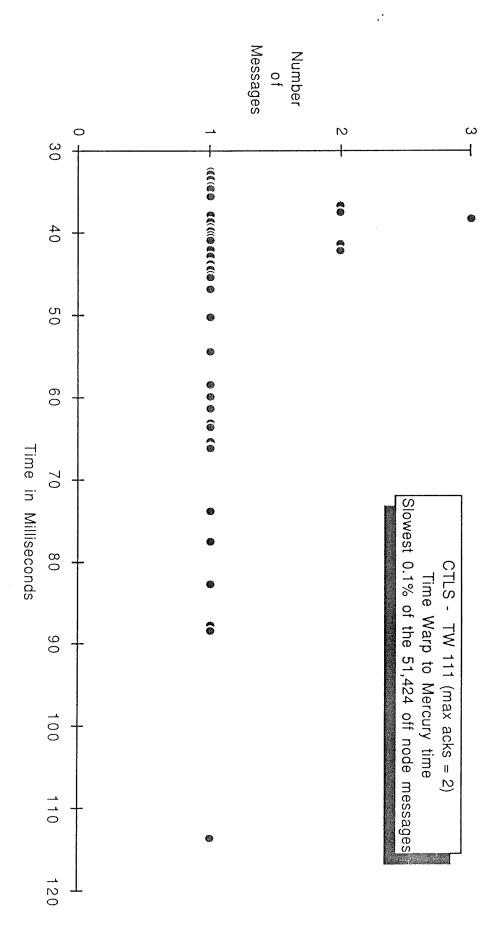


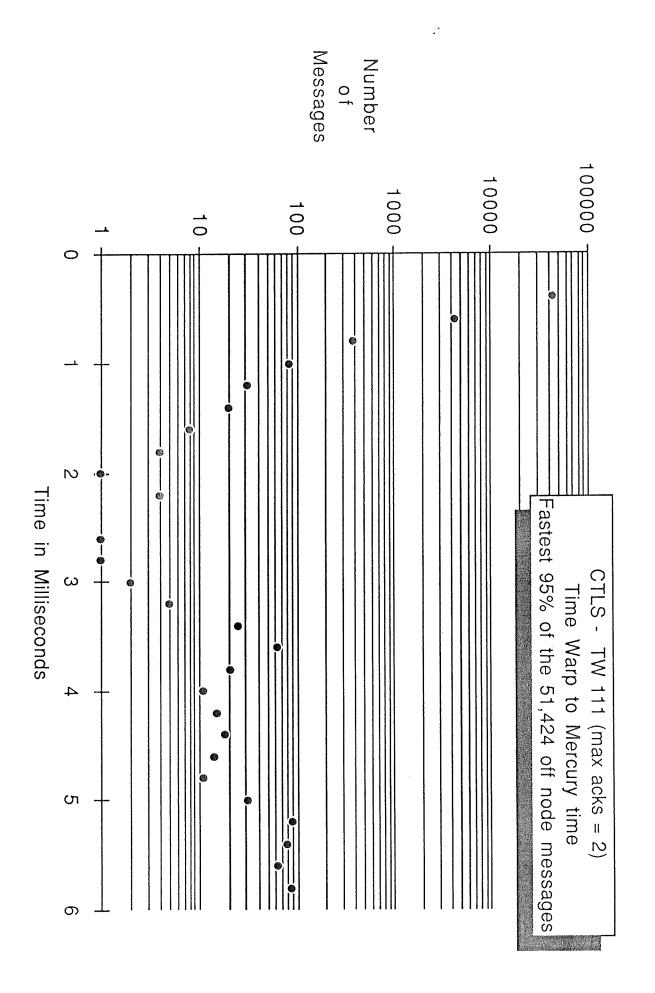
Appendix B Time Warp to Mercury Times

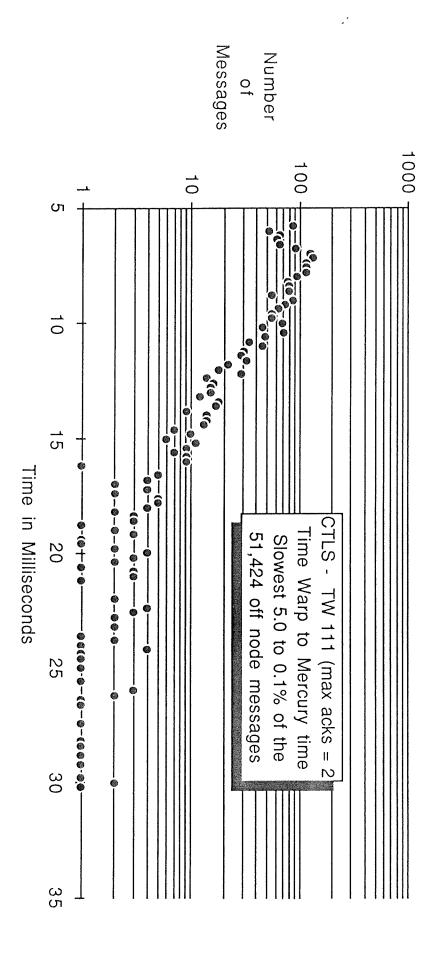
- B-1: Fastest 95%, acks=2, graph (see B-4).
- B-2: The 95% 99.9% range, acks=2, graph (see B-5).
- B-3: Slowest 0.1%, acks=2, graph (see B-6).
- B-4: Fastest 95%, acks=2, semi-log graph (see B-1).
- B-5: The 95% 99.9% range, acks=2, semi-log graph (see B-2).
- B-6: Slowest 0.1%, acks=2, semi-log graph (see B-3).
- B-7: 100%, acks=20, graph (see B-8).
- B-8: 100%, acks=20, semi-log graph (see B-7).
- B-9: Fastest 99.96%, acks=20, graph (see B-10).
- B-10: Fastest 99.96%, acks=20, semi-graph (see B-9).

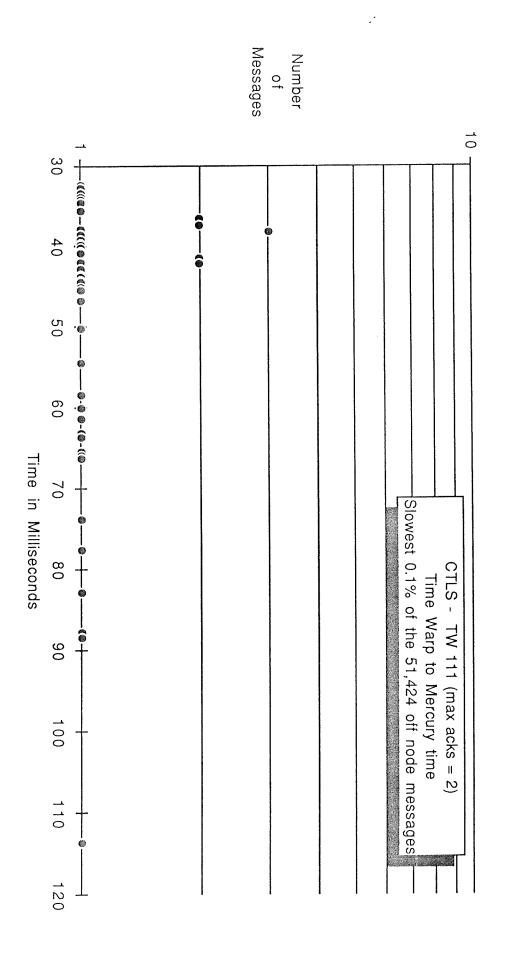


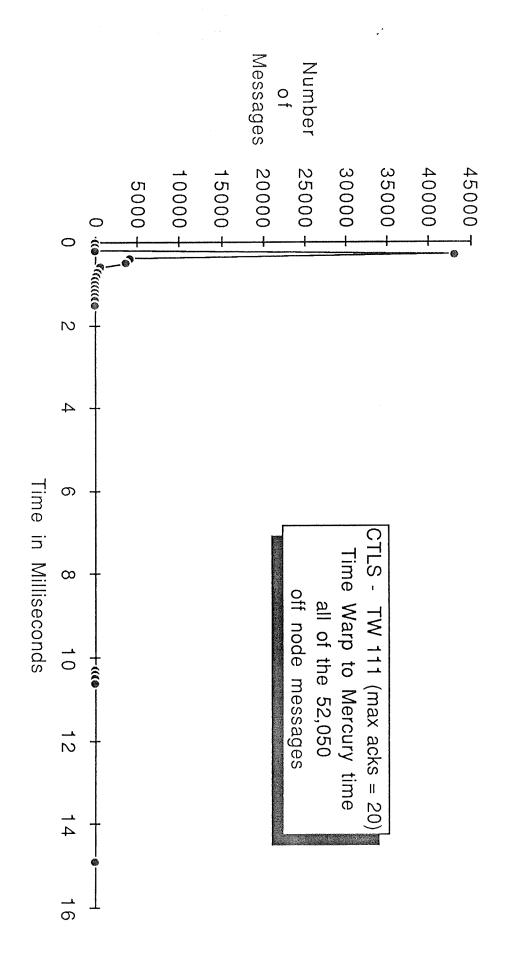


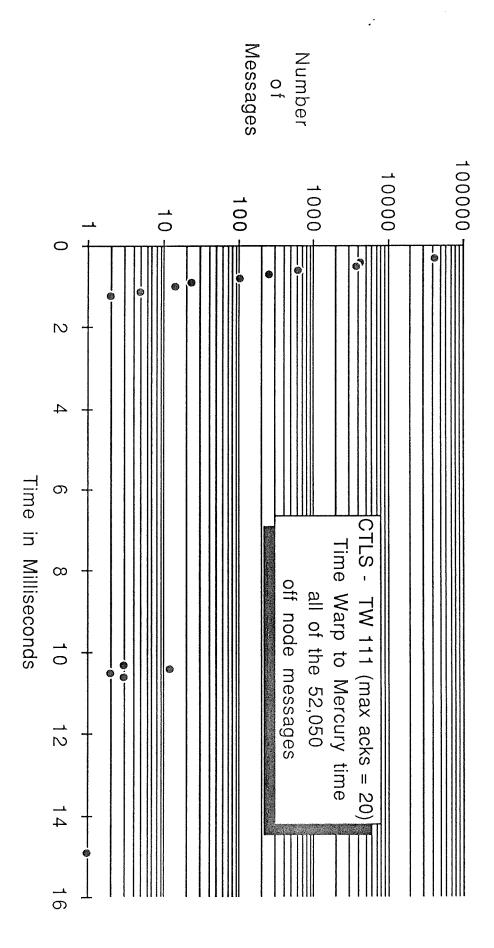


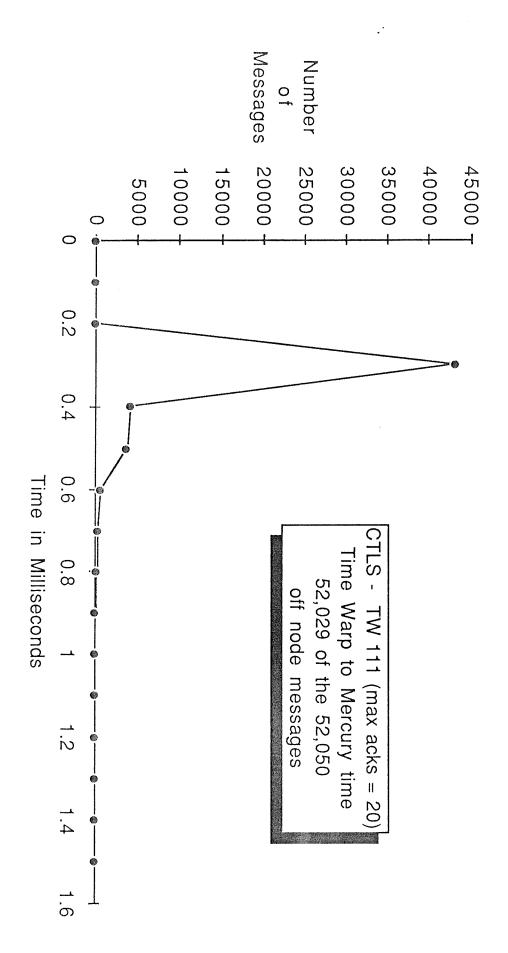


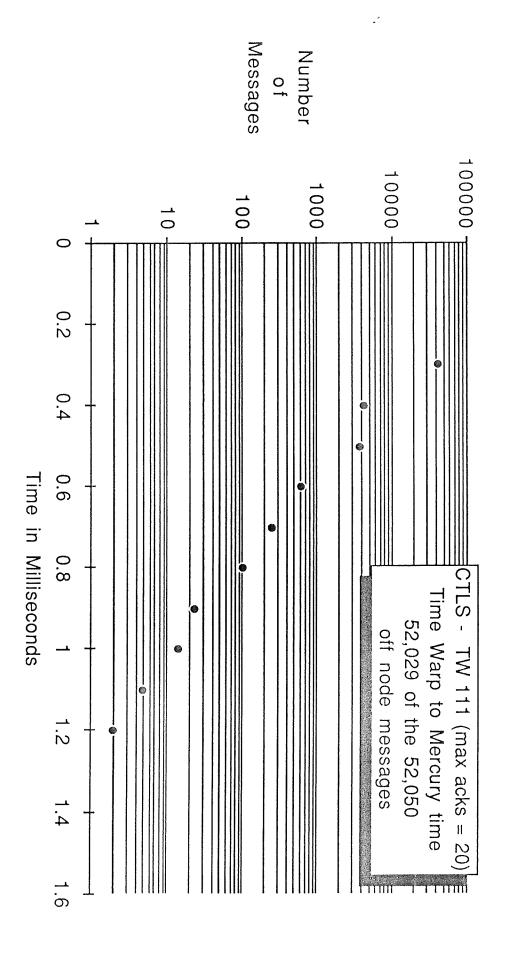






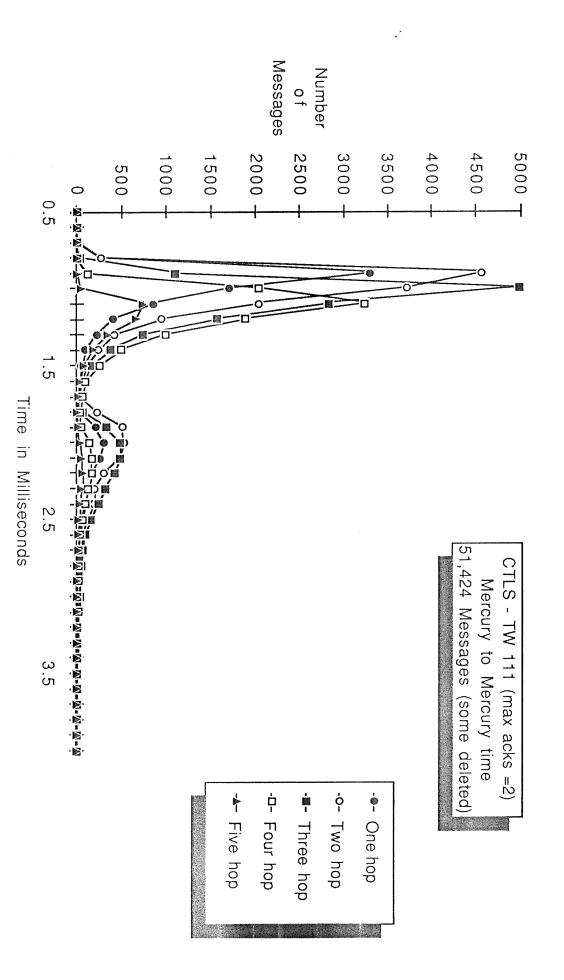


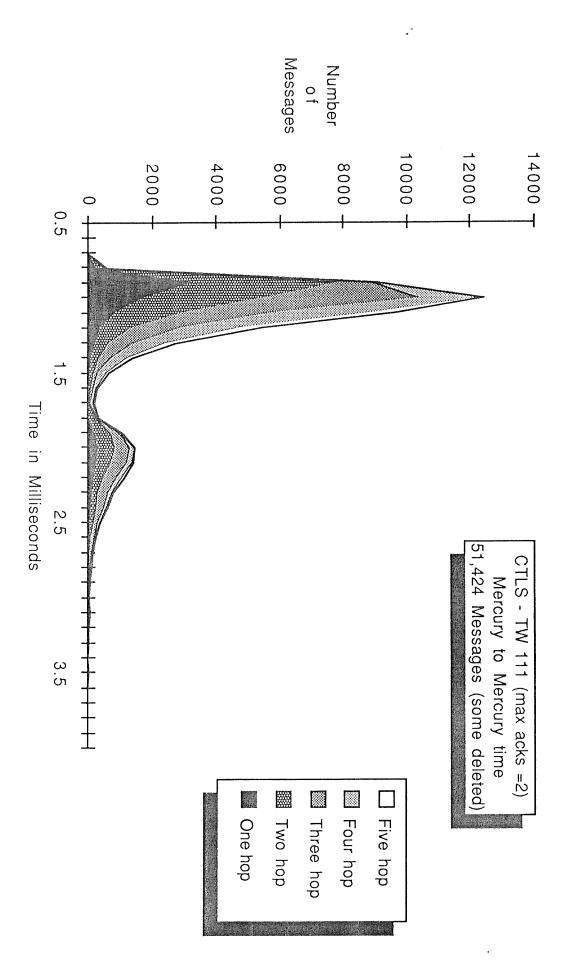


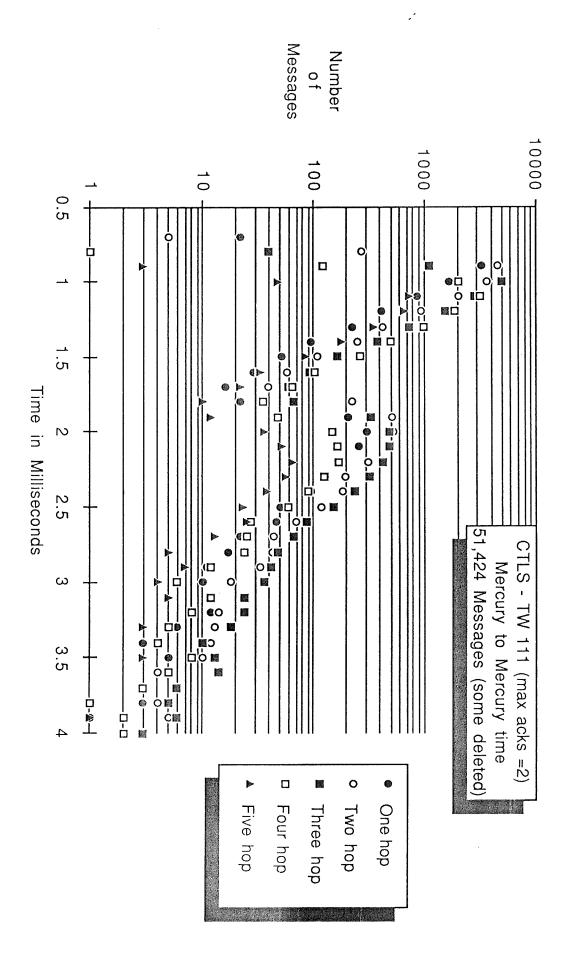


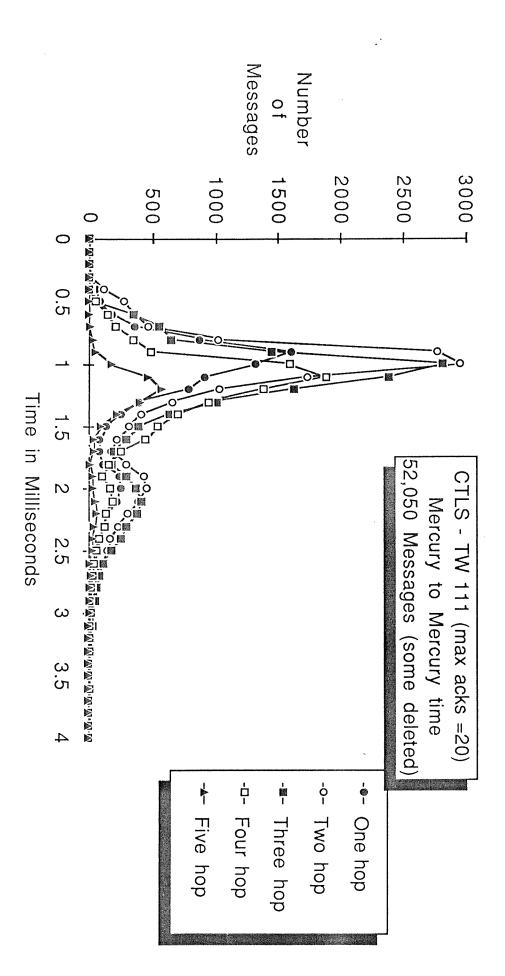
Appendix C Mercury to Mercury Times

- C-1: Fastest 99.66%, acks=2, by hops, graph (see C-2, C-3).
- C-2: Fastest 99.66%, acks=2, by hops, area graph (see C-1, C-3).
- C-3: Fastest 99.66%, acks=2, by hops, semi-log graph (see C-1, C-2).
- C-4: Fastest 99.64%, acks=20, by hops, graph (see C-5, C-6).
- C-5: Fastest 99.64%, acks=20, by hops, area graph (see C-4, C-6).
- C-6: Fastest 99.64%, acks=20, by hops, semi-log graph (see C-5, C-6).
- C-7: Slowest 0.36%, acks=20, by hops graph.
- C-8: Fastest 3 hop, acks=2, by packets, area graph (see C-9).
- C-9: Fastest 3 hop, acks=2, by packets, semi-log graph (see C-8).

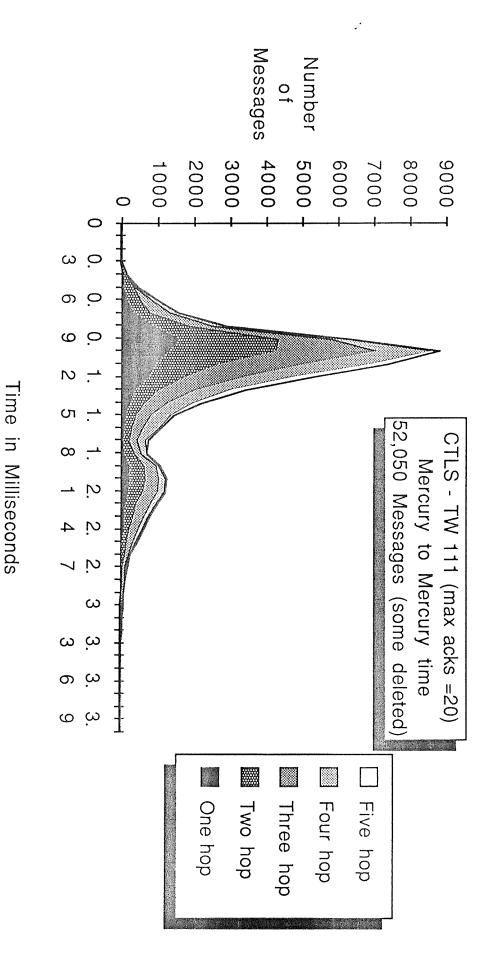


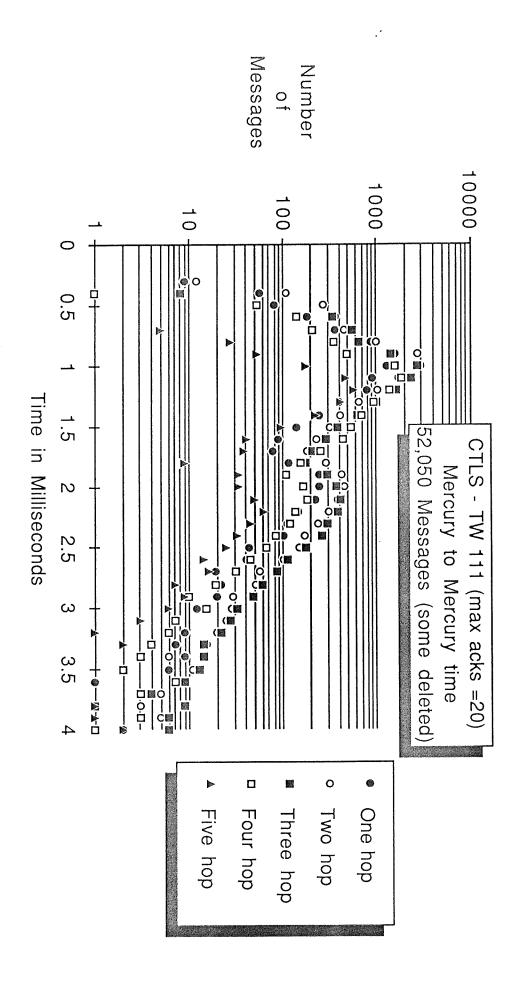


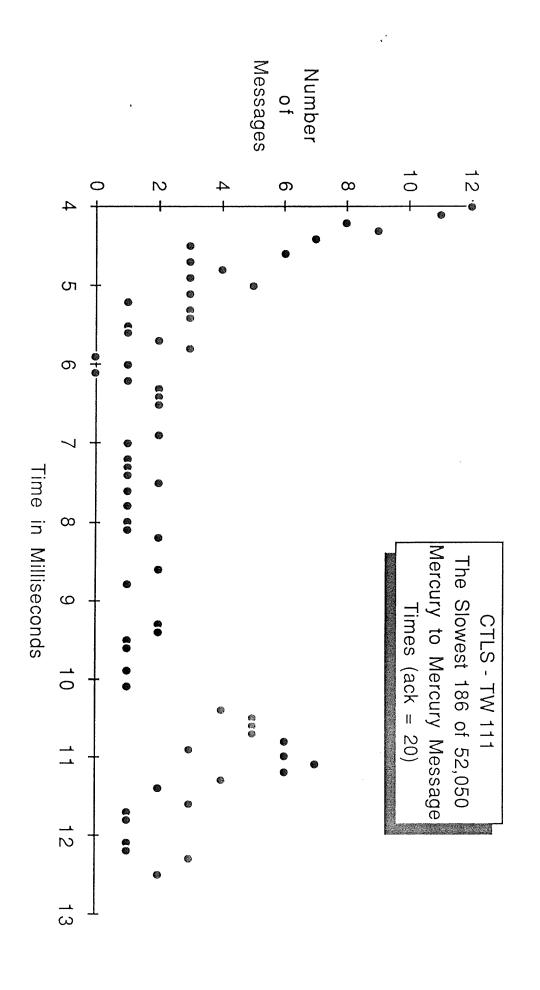


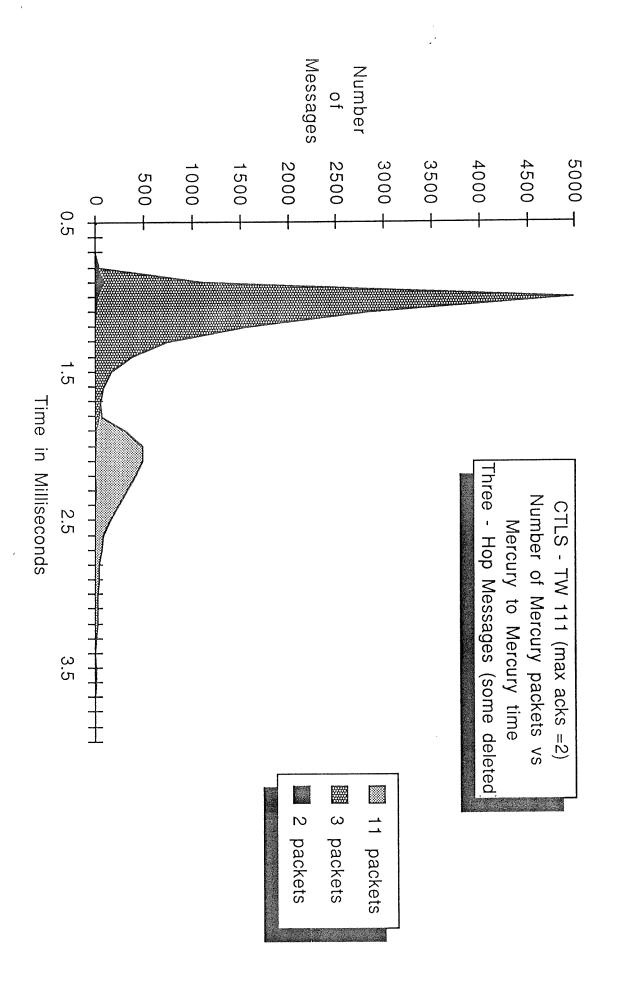


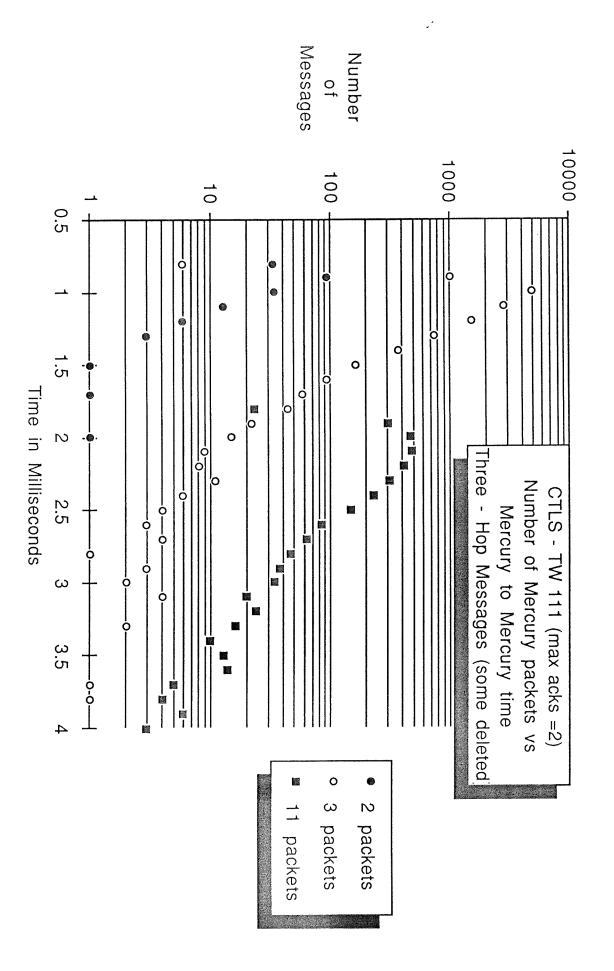
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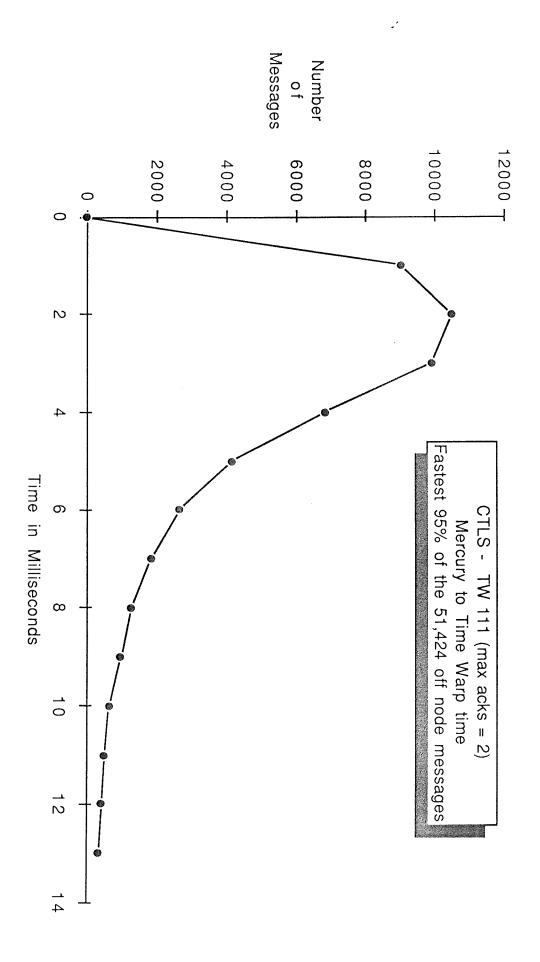


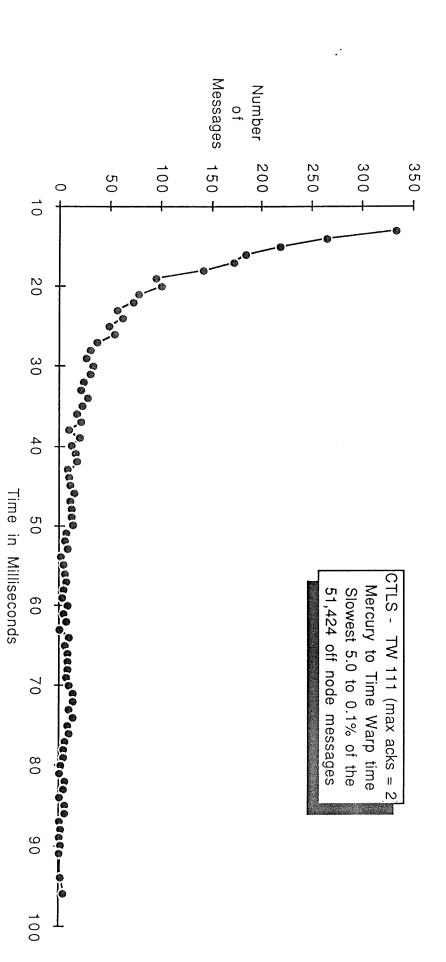


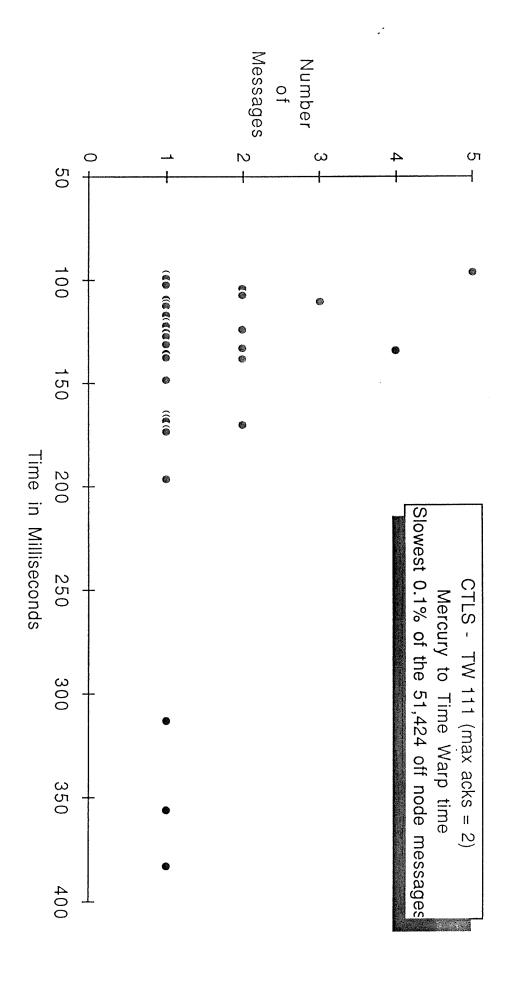


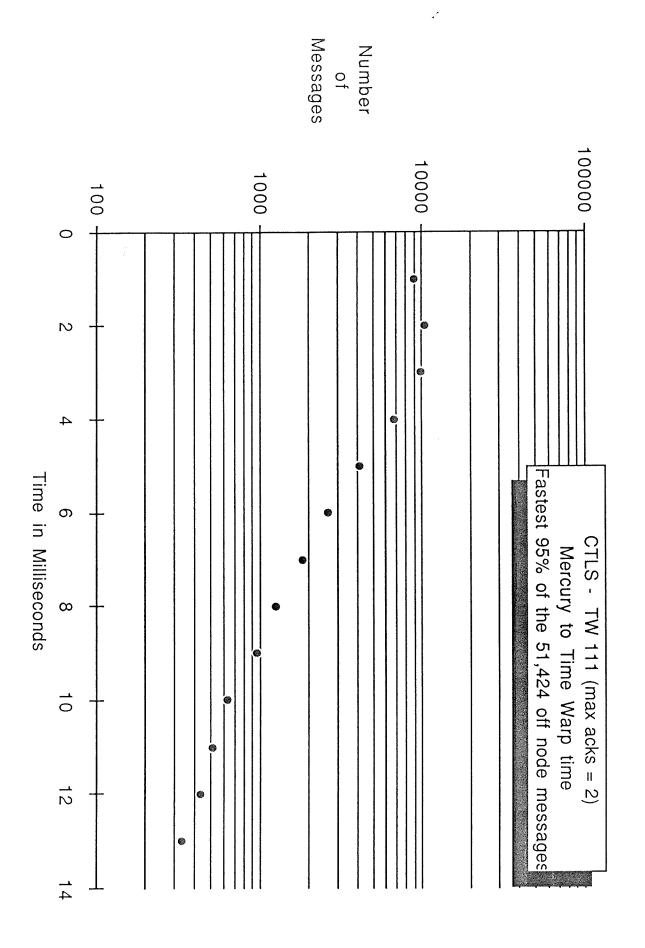
Appendix D Mercury to Time Warp Times

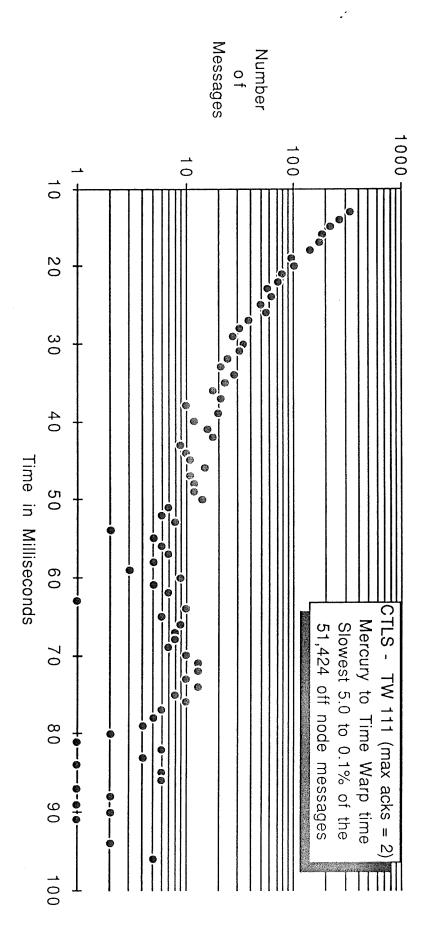
- D-1: Fastest 95%, acks=2, graph (see D-4).
- D-2: The 95% 99.9% range, acks=2, graph (see D-5).
- D-3: Slowest 0.1%, acks=2, graph (see D-6).
- D-4: Fastest 95%, acks=2, semi-log graph (see D-1).
- D-5: The 95% 99.9% range, acks=2, semi-log graph (see D-2).
- D-6: Slowest 0.1%, acks=2, semi-log graph (see D-3).

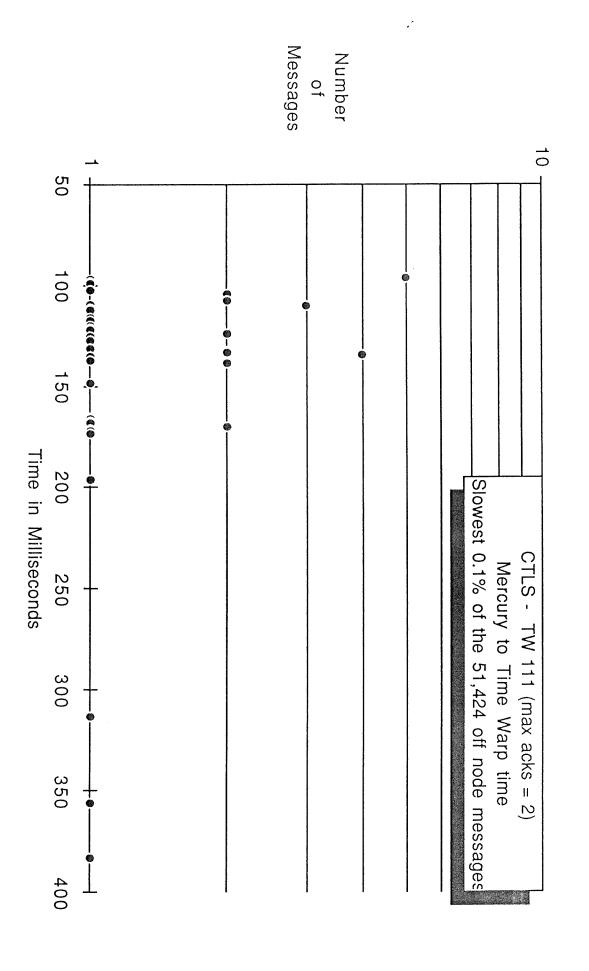






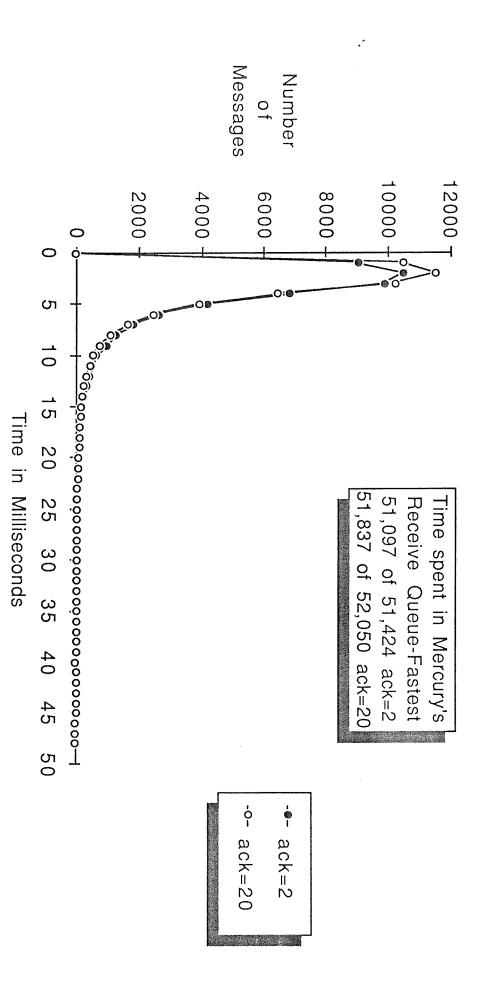


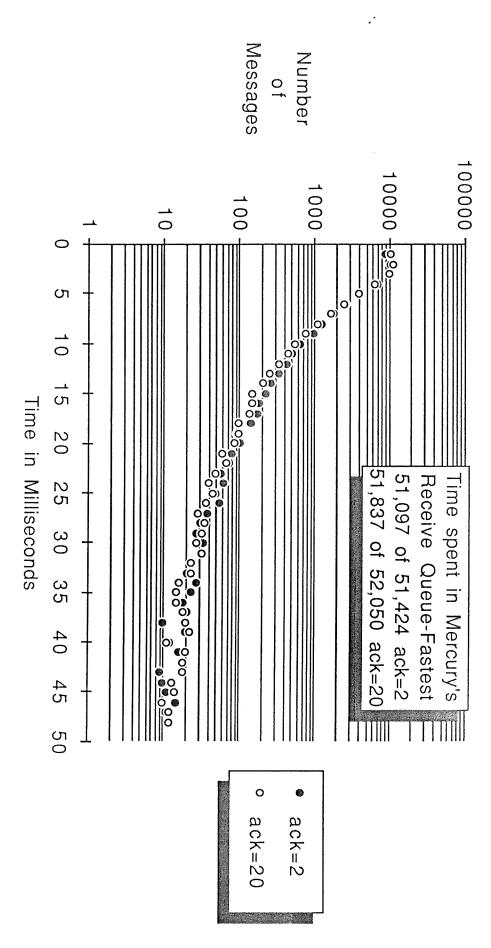




Appendix E Queueing Delays Time in Mercury's Receive Queue Length of Mercury's Receive Queue

- E-1: Fastest times, acks=2 vs acks=20, graph (see E-2).
- E-2: Fastest times, acks=2 vs acks=20, semi-log graph (see E-1).
- E-3: Slowest times, acks=2 vs acks=20, graph (see E-4).
- E-4: Slowest times, acks=2 vs acks=20, semi-log graph (see E-3).
- E-5: Length, acks=2, graph (see E-6).
- E-6: Length, acks=2, semi-log graph (see E-5).
- E-7: Length, acks=2 vs acks=20, semi-log graph.
- E-8: On node length, acks=20, semi-log graph.
- E-9: On node length, acks=2, semi-log graph.





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