

## MAD 3105 — Discrete Math 2

Section 2, Fall 1993.

Instructor: Bellenot.

The good doctor's Office: 002-B Love, Office Hours: MW 1:30–2:20 or by appointment.

Eligibility: A grade of C- or better in Discrete 1 (MAD 3104)

Text: Mott, Kandel, Baker; 2nd Edition.

Coverage: Some or all of chapters 2, 3, 4, 5 and 7.

Final: At 7:30 – 9:30 am Thursday Dec 16, 1993.

Tests: (3) Tentatively at Sept 22, Oct 27(Nov 5?) and Dec 1. No Makeup tests.

Quizzes: Every Wednesday (except test days). No Makeup quizzes.

Grades: 90% A, 80%B, 70%C, 60%D.

Relative Weights  $F = 2T$  and  $T = Q$  (F is  $1/3$ , each T is  $1/6$  and Q is  $1/6$ ).

Homework and Attendance are required.

Fair Warning: The good doctor teaches this course differently than other instructors. The order of material will be different and so will be the emphasis.

Show **ALL** work for credit; be neat; and use only **ONE** side of each page of paper.

- Write the general solution to the homogenous linear recurrence relation if the characteristic polynomial is
  - $x^2 - 5x + 6$ .
  - $x^2 + 4x + 4$ .
  - $x^2 - 2x + 1$ .
  - $x^2 - x - 1$ .
- Give the *correct guesses* to the form of the particular solution of the linear non-homogenous recurrence relation which has the given forcing function. The characteristic polynomial is known to have roots: -5, 1, 1, 3, 3, 3, 3, 5 and 6.
  - $145 \cdot 2^n$ .
  - $(-n^2 + 7n - 5)4^n$ .
  - $65 \cdot 5^n$ .
  - $n^2 + 7$ .
- Show that the substitution  $b_n = \log_2(a_n)$  transforms the non-linear recurrence relation  $a_n = (a_{n-1})^2(a_{n-2})^3; a_0 = 4, a_1 = 8$ , into a linear recurrence relation. Do **NOT** solve these recurrence relations. Do translate the initial conditions.
- Simplify** the following to the form  $n^\alpha$ .
  - $7^{\log_3 n}$ .
  - $(7^\pi)^{\log_5 n}$ .
- Show  $a_n = n \log_2(n)$  is a solution to  $a_n = 2a_{n/2} + n, a_1 = 0$ ; for  $n = 2^k \geq 1$ .
- Solve:  $a_n - a_{n-1} - 6a_{n-2} = 0; a_0 = 2, a_1 = 6$ .
- Solve:  $a_n - a_{n-1} = n + 1; a_0 = 7$ .
- Solve by **GENERATING FUNCTIONS**:  $a_n = 5a_{n-1} + 3^n, n \geq 1; a_0 = 2$ .

Show **ALL** work for credit; be neat; and use only **ONE** side of each page of paper.

1. Write the general solution to the homogenous linear recurrence relation if the characteristic polynomial is

A.  $x^2 - 5x + 6$ .

Answer:  $A2^n + B3^n$

B.  $x^2 + 4x + 4$ .

Answer:  $A(-2)^n + Bn(-2)^n$

C.  $x^2 - 2x + 1$ .

Answer:  $An + B$

D.  $x^2 - x - 1$ .

Answer:  $A((1 - \sqrt{5})/2)^n + B((1 + \sqrt{5})/2)^n$

2. Give the *correct guesses* to the form of the particular solution of the linear non-homogenous recurrence relation which has the given forcing function. The characteristic polynomial is known to have roots: -5, 1, 1, 3, 3, 3, 3, 5 and 6.

A.  $145 \cdot 2^n$ .

Answer:  $A2^n$

B.  $(-n^2 + 7n - 5)4^n$ .

Answer:  $(An^2 + Bn + C)4^n$

C.  $65 \cdot 5^n$ .

Answer:  $nA5^n$

D.  $n^2 + 7$ .

Answer:  $n^2(An^2 + Bn + C)$

3. Show that the substitution  $b_n = \log_2(a_n)$  transforms the non-linear  $a_n = (a_{n-1})^2(a_{n-2})^3$ ;  $a_0 = 4, a_1 = 8$ , into a linear recurrence relation. Do **NOT** solve these recurrence relations. Do translate the initial conditions.

Answer: First the initial conditions:  $b_0 = \log_2 a_0 = \log_2 4 = 2$  and  $b_1 = \log_2 a_1 = \log_2 8 = 3$ . Second, taking  $\log_2$  of both sides of the original recurrence relation yields  $\log_2 a_n = 2 \log_2 a_{n-1} + 3 \log_2 a_{n-2}$  or  $b_n = 2b_{n-1} + 3b_{n-2}$  which is a linear recurrence relation.

4. Simplify the following to the form  $n^\alpha$ .

A.  $7^{\log_3 n}$ .

Answer:  $n^{\log_3 7}$

B.  $(7^\pi)^{\log_5 n}$ .

Answer:  $n^{\pi \log_5 7}$

5. Show  $a_n = n \log_2(n)$  is a solution to  $a_n = 2a_{n/2} + n, a_1 = 0$ ; for  $n = 2^k \geq 1$ .

Answer: First the initial condition  $a_1 = 1 \log_2 1 = 1 \cdot 0 = 0$ , which is correct. Second substituting we get see  $RHS = 2a_{n/2} + n = 2(n/2) \log_2(n/2) + n = n(\log_2 n - \log_2 2) + n = n(\log_2 n - 1) + n = n \log_2 n - n + n = n \log_2 n = a_n = LHS$ .

6. Solve:  $a_n - a_{n-1} - 6a_{n-2} = 0; a_0 = 2, a_1 = 6$ .

Answer: The general solution is  $a_n = A(-2)^n + B3^n$ . Using the initial conditions we get  $2 = a_0 = A + B$  and  $6 = a_1 = -2A + 3B$ . Adding twice the first equation to the second yields  $10 = 5B$  so  $B = 2$  and  $A = 0$ . The answer is  $a_n = 2 \cdot 3^n$ .

7. Solve:  $a_n - a_{n-1} = n + 1; a_0 = 7$ .

Answer: First the particular solution has form  $a_n = (An + B)n = An^2 + Bn$ . Substituting  $An^2 + Bn - A(n-1)^2 - B(n-1) = n^2(A-A) + n(B+2A-B) + 1(-A+B)$  so we solve  $2A = 1$  and  $B - A = 1$  so  $A = 1/2, B = 3/2$ . Thus the general solution is  $a_n = C + n^2/2 + 3n/2$  and when  $n = 0, 7 = a_0 = C + 0 + 0$ , so  $C = 7$  and  $a_n = 7 + 3n/2 + n^2/2$ .

8. Solve by **GENERATING FUNCTIONS**:  $a_n = 5a_{n-1} + 3^n, n \geq 1; a_0 = 2$ .

Answer:  $A(X) = \sum_{n=0}^{\infty} a_n X^n = a_0 + \sum_{n=1}^{\infty} a_n X^n = 2 + \sum_{n=1}^{\infty} (5a_{n-1} + 3^n) X^n = 2 + 5X \sum_{n=1}^{\infty} a_{n-1} X^{n-1} + 3X \sum_{n=1}^{\infty} 3^{n-1} X^{n-1} = 2 + 5X \sum_{n=0}^{\infty} a_n X^n + 3X \sum_{n=0}^{\infty} 3^n X^n = 2 + 5XA(X) + \frac{3X}{1-3X}$  Solving for  $A(X)$  yields

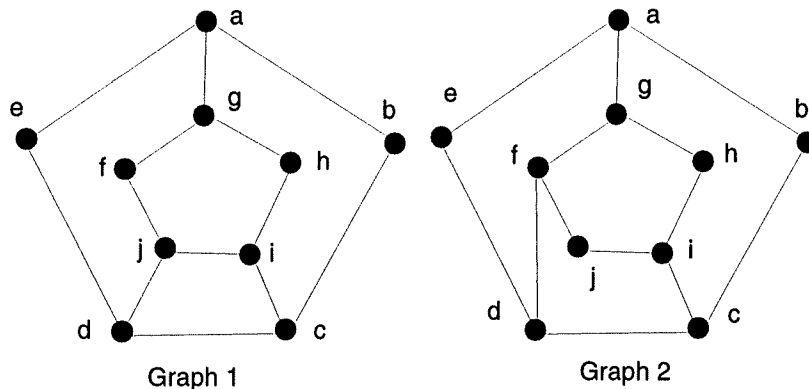
$$A(X)(1 - 5X) = \frac{2 - 6X}{1 - 3X} + \frac{3X}{1 - 3X} \cdot A(x) = \frac{2 - 3X}{(1 - 5X)(1 - 3X)} = \frac{B}{1 - 5X} + \frac{C}{1 - 3X}$$

$$B = \frac{2 - 3(1/5)}{1 - 3(1/5)} = \frac{7/10}{2/5} = \frac{7}{4}, C = \frac{2 - 3(1/3)}{1 - 5(1/3)} = \frac{1}{-2/3} = \frac{-3}{2}$$

Thus  $A(x) = (7/4) \sum_{n=0}^{\infty} 5^n + (-3/2) \sum_{n=0}^{\infty} 3^n$  and hence  $a_n = (7/4)5^n + (-3/2)3^n$ .

Show **ALL** work for credit; be neat; and use only **ONE** side of each page of paper.

1. What big  $O$  (i.e.  $O(f(n))$ ) best describes the run time of the following algorithms:
- Sorting  $n$  items by Merge Sort.
  - Searching in a sorted array of  $n$  items by Binary Search.
  - Sorting  $n$  items by Interchange Sort.
- DE. Where the run time is given by  $a_n = 4a_{n/2}; a_1 = 1$ . Simplify.

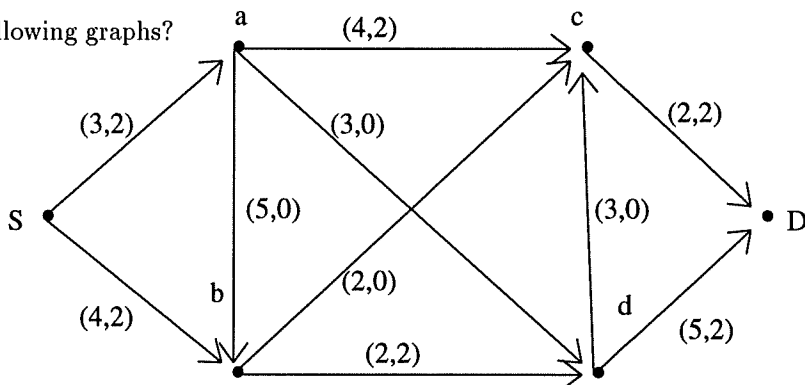


2. For Graph 1 above:
- Either construct an Euler circuit or prove none exists.
  - Either construct a Hamiltonian cycle or prove none exists.
  - Draw the dual graph.
3. Disorder and more big  $O$ .
- Prove  $O(4n^2 - 5 + 13n^3) = O(100n^3 - n + n^2 - 3)$ .
  - Arrange in increasing order:  
 $O(\log n), O(n!), O(n^6), O(\sqrt{n}), O(3^n), O(n), O(n\sqrt{n}), O(n^n), O(n \log n), O(2^n)$ .
  - Compute the disorder of 2, 7, 3, 1, 4, 5, 9, 8, 6.
4. Show that Graph 2 above does not have a Hamiltonian cycle:
- Using Grinberg's Theorem.
  - Without using Grinberg's Theorem.
5. For the relation  $R$  on the set of integers defined by  $xRy \iff 0 < |x - y| < 3$ . Either prove  $R$  has the given property or give a counter-example to show it does not have that property.
- Reflexive.
  - Symmetric.
  - Anti-symmetric.
  - Transistive.
6. Mod  $m$  problems.
- Find  $0 \leq x < 31$  so that  $7x \equiv 1 \pmod{31}$
  - Find  $0 \leq x < 5$  so that  $3^{101} \equiv x \pmod{5}$
  - What is the remainder if  $\sum_{i=1}^{100} i^2$  is divided by 5.
7. Prove the following relation  $R$  is an equivalence relation and describe the equivalence class of  $(3,4)$ . The relation  $R$  is defined on the points of the plane by  $(a,b)R(c,d) \iff a^2 + b^2 = c^2 + d^2$ .
8. Use Euler's formula and  $2|E| \geq 3|R|$  to show  $|E| \leq 3|V| - 6$ . Then show a plane connected (simple) graph with strictly fewer than 30 edges has a vertex of degree 4 or less.

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1. What is the chromatic number of the following graphs?

- A.  $K_n$
- B.  $\overline{K_n}$
- C.  $C_n, n \geq 4, n$  even.
- D.  $C_n, n \geq 3, n$  odd.
- E. A non-trivial tree.



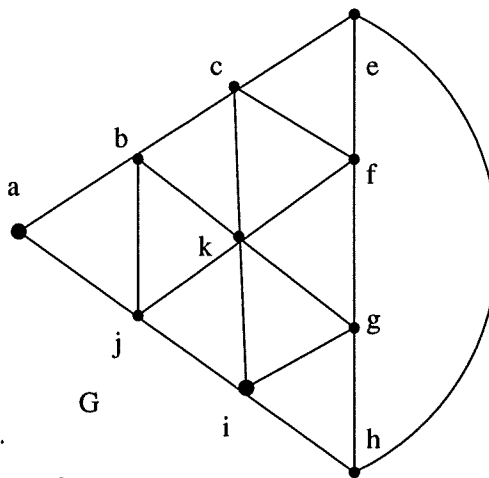
2. For the transport network above use our labeling algorithm to find a maximal flow, its value, and a minimal cut.

3. Network True or False.

- A. If  $F(X, \overline{X}) - F(\overline{X}, X) = k(X, \overline{X})$ , then  $F$  is a maximal flow.
- B. If  $F(\overline{X}, X) = 0$  and  $(X, \overline{X})$  is minimal cut, then  $F$  is a maximal flow.
- C. A cut can be minimal for one flow, and not minimal for another flow.
- D. If  $ab$  is an unsaturated edge with non-zero flow and  $(X, \overline{X})$  is a minimal cut, then either both vertices are in  $X$  or both vertices are in  $\overline{X}$ .
- E. If  $|F| = 0$ , then every edge has zero flow.

4. Give network counterexamples to each statement below:

- A. A transport network with a unique maximal flow has a unique minimal cut.
- B. A transport network with a unique minimal cut has a unique maximal flow.
- C. If  $F$  is a flow and  $(X, \overline{X})$  is a cut so that  $F(\overline{X}, X) = 0$ , then  $(X, \overline{X})$  is a minimal cut.



5. For  $G$  above

- A. Find an Euler circuit or prove none exists.
- B. Find an Euler path or prove none exists.
- C. Find an Hamiltonian cycle or prove none exists.

6. Find  $\chi(G)$  for  $G$  above. Prove  $G$  has this chromatic number.

7. Using  $2|E| \leq 3|R|$  and Euler's formula show plane connected (simple) graphs satisfies  $|E| \leq 3|V| - 6$ . Show plane connected (simple) graphs have a vertex of degree 5 or less.

8. There are exactly two 3-regular graphs with 6 vertices,  $K_{3,3}$  and the "triangular prism". For both of these graphs, either give a plane drawing of the graph, or carefully prove the graph is non-planar.